

THE CHOICE OF OPTIMAL DECISIONS ÎN UNCERTITUDE SITUATIONS

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Abstract: *The Electre method is a classical algorithm for the decidence of a suitable choice în the process of launching on market of some products. In this paper I shall give a variation of the last part of the algorithm în the direction of simplifying the finally computations.*

In many practical situations, it exists many informations refered to the actions which can be developed, but for each action choice it exists many possibilities.

The Electre method gives us a way which will make out between different possibilities.

Let therefore n choices V_1, V_2, \dots, V_n for a decident. We have also m criterions C_1, C_2, \dots, C_m who have, each of them, a weighty coefficient (by rule subjective assigned) k_1, k_2, \dots, k_m .

For each pair (V_i, C_j) we have a numerical value (if it is a qualitative estimation we shall convert în hierarchy numbers).

The algorithm will determine the best choice of action.

Step I

We settle, for the beginning the method nature: maximisation or minimisation. We shall add two lines under the table, on which will compute the maximum and the minimum of all numbers on columns.

Step II

We shall determine the utilities U_{ij} corresponding to all pairs (V_i, C_j) :

- ◆ for the maximisation problem: $U_{ij} = \frac{v_{ij} - \min_{k=1,\dots,n} v_{kj}}{\max_{k=1,\dots,n} v_{kj} - \min_{k=1,\dots,n} v_{kj}}$;
- ◆ for the minimisation problem: $U_{ij} = \frac{\max_{k=1,\dots,n} v_{kj} - v_{ij}}{\max_{k=1,\dots,n} v_{kj} - \min_{k=1,\dots,n} v_{kj}}$.

and after we shall draw the table with these results.

Step III

We shall compute the correspondence indicators: $c(V_i, V_j) = \frac{\sum_{\substack{p=1,\dots,m \\ U_{ip} \geq U_{jp}}} k_p}{\sum_{r=1}^m k_r}$ and after the disunity indicators: $d(V_i, V_j) = \max_{p=1,\dots,m} (U_{jp} - U_{ip}, 0)$.

Step IV

We shall draw a new table for each pair (V_i, V_j) where on each cell we shall write on the left the correspondence indicator and on the right the disunity indicator.

Step V

We shall establish two values p and q (with a significance of complementary probabilities) such that $p, q \in (0, 1)$ and $p + q = 1$ which will extent the admissible limits for correspondence, respectively disunity. We shall say that V_i is better than V_j if

$$\begin{cases} c(V_i, V_j) \geq p \\ d(V_i, V_j) \leq q \end{cases}$$

Let now the matrice $G = (g_{ij}) \in M_n(\mathbb{R})$ such that: $g_{ij}=1$ if V_i is better than V_j and 0 in the other situations. If it exists a row i with all elements equal with 1 it follows that V_i is better than all others therefore it will be preferred. If it not exists such a V_i we shall diminish the value of p (and of course increase q) till we shall obtain the better choice.

In what follows I shall present another way to choice p and q .

The system: $\begin{cases} c(V_i, V_j) \geq p \\ d(V_i, V_j) \leq q \end{cases}$ can be write as: $\begin{cases} c(V_i, V_j) \geq p \\ d(V_i, V_j) \leq 1-p \end{cases}$ therefore:

$$p \leq \min(c(V_i, V_j), 1-d(V_i, V_j))$$

We shall compute $\min_{j=1,n} c(V_i, V_j)$ and $\min_{j=1,n} (1-d(V_i, V_j))$ from where:

$$p \leq \min\left(\min_{j=1,n} c(V_i, V_j), 1 - \max_{j=1,n} d(V_i, V_j)\right)$$

In the last table, we shall add therefore three columns, on which we compute $\min_{j=1,n} c(V_i, V_j)$, $1 - \max_{j=1,n} d(V_i, V_j)$ and the minimum of these two values. The better choice will be those who give the maximum value on this last column.

Example

Let the following problem (of maximisation):

Criterion Choice	C ₁ (k ₁ =0,4)	C ₂ (k ₂ =0,4)	C ₃ (k ₂ =0,2)
V ₁	1000	0	50
V ₂	800	1	56
V ₃	600	2	60
V ₄	700	1	54
V ₅	500	2	58

Criterion Choice	C ₁	C ₂	C ₃
V ₁	1000	0	50
V ₂	800	1	56
V ₃	600	2	60
V ₄	700	1	54
V ₅	500	2	58
min	500	0	50
max	1000	2	60
max-min	500	2	10

The utilities table is:

Criterion Choice	C ₁ (k ₁ =0,4)	C ₂ (k ₂ =0,4)	C ₃ (k ₃ =0,2)
V ₁	1	0	0
V ₂	0,6	0,5	0,6
V ₃	0,2	1	1
V ₄	0,4	0,5	0,4
V ₅	0	1	0,8

And the table of correspondence indicators and the disunity indicators:

	V ₁		V ₂		V ₃		V ₄		V ₅	
	V ₁	V ₂	V ₂	V ₃	V ₃	V ₄	V ₄	V ₅	V ₅	V ₅
V ₁	1	0	0,4	0,6	0,4	1	0,4	0,5	0,4	1
V ₂	0,6	0,4	1	0	0,4	0,5	1	0	0,4	0,5

V₃	0,6	0,8	0,6	0,4	1	0	0,6	0,2	1	0
V₄	0,6	0,6	0,4	0,2	0,4	0,6	1	0	0,4	0,5
V₅	0,6	1	0,6	0,6	0,4	0,2	0,6	0,4	1	0

If we try for p from 1 back to 0 we shall obtain that, for the first time, we shall have for p=0,4 and q=0,6:

$$G = \begin{pmatrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \end{pmatrix}$$

therefore each of V₂ and V₄ is better than the others.

I shall suggest at this last step the following:

	V₁		V₂		V₃		V₄		V₅		$\min_{j=1,n} c(V_i, V_j)$	$\max_{j=1,n} d(V_i, V_j)$	min
V₁	1	0	0,4	0,6	0,4	1	0,4	0,5	0,4	1	0,4	1	0
V₂	0,6	0,4	1	0	0,4	0,5	1	0	0,4	0,5	0,4	0,5	0,4
V₃	0,6	0,8	0,6	0,4	1	0	0,6	0,2	1	0	0,6	0,8	0,2
V₄	0,6	0,6	0,4	0,2	0,4	0,6	1	0	0,4	0,5	0,4	0,6	0,4
V₅	0,6	1	0,6	0,6	0,4	0,2	0,6	0,4	1	0	0,4	1	0

therefore each of V₂ and V₄ is better than the others, but without consecutive tests.